

Regret-Minimizing Project Choice

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Motivating example I

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- The department proposes one candidate to the dean, who decides whether to make an offer or not

Motivating example II

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- The employee proposes a project to the headquarters, which decides whether to approve or not

Motivating example III

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- A division has a chance to choose an office building
- It learns that the available locations are $\{L, M, N, P\}$
- The division proposes a short list for the headquarters to evaluate and choose

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 - an agent learns which projects are available and proposes project(s)
 - a principal evaluates the proposed project(s) and makes the choice
- Proposing bias: the agent has a tendency to propose his favorite project and hide his less preferred ones

Questions

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What is the strategic role of multiproject proposals?

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What is the strategic role of multiproject proposals?

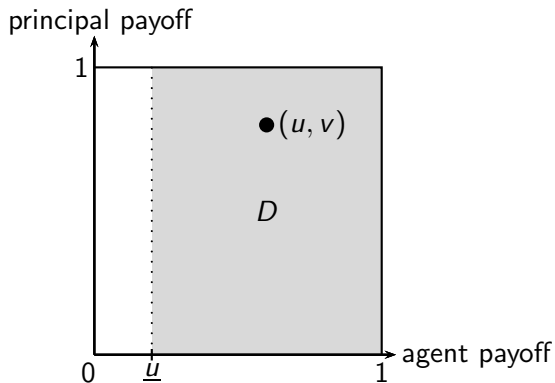
Do several projects within the proposal have a chance of being chosen?

Roadmap

- Model
- Single-project environment
- Multiproject environment
- Discussion

Projects and payoffs

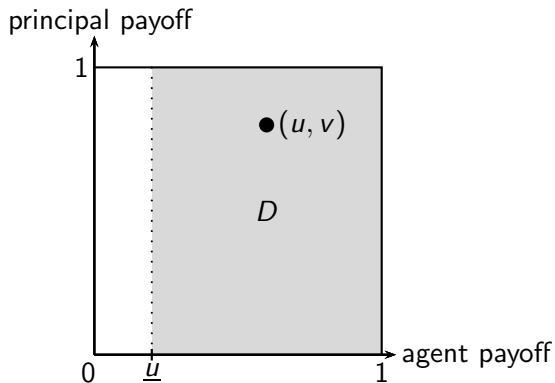
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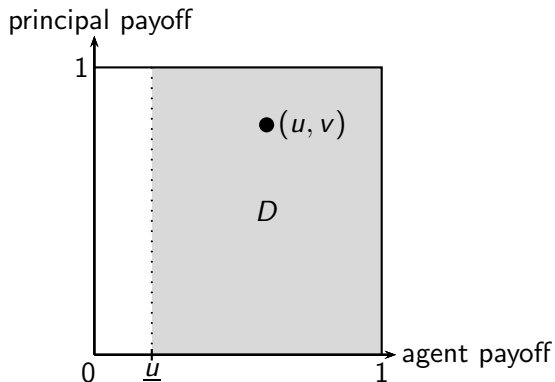
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For the talk $\underline{v} = 0$, $\underline{u} \in [0, 1]$



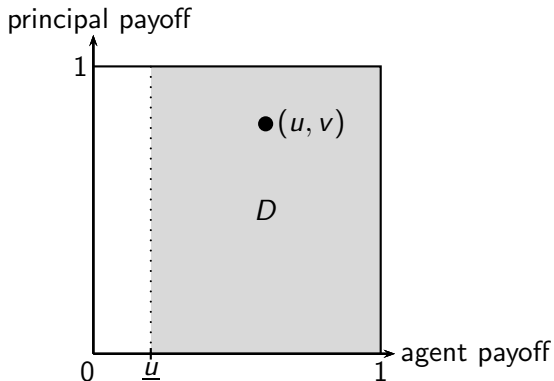
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- If project (u, v) is chosen: the agent gets u ; the principal gets v
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multi-project environment: $K = \infty$

intermediate environment: $2 \leq K < \infty$

Mechanism

- A mechanism ρ attaches to each proposal P a subprobability measure $\rho(\cdot|P)$ over P :

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- The principal commits to a mechanism ρ

The principal's problem

Given mechanism ρ and the agent's type A :

- the agent chooses a proposal P to maximize his expected payoff:

$$P \in \operatorname{argmax}_{P \subseteq A, |P| \leq K} \sum_{(u,v) \in P} \rho((u,v)|P)u$$

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- the principal's regret is

$$\operatorname{RGRT}(\rho, A) = \max_{(u,v) \in A} v - \sum_{(u,v) \in P} \rho((u,v)|P)v$$

The principal's problem

- The principal's worst-case regret under mechanism ρ is

$$\text{WCR}(\rho) = \sup_{A \subseteq D, |A| < \infty} \text{RGRT}(\rho, A),$$

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- The principal chooses ρ that minimizes his worst-case regret

Related literature

- **Project choice:** Armstrong and Vickers (2010), Nocke and Whinston (2013)
- **Mechanism design with worst-case regret:** Hurwicz and Shapiro (1978), Bergemann and Schlag (2008, 2011), Manski (2011), Renou and Schlag (2011), Beviá and Corchón (2019), Malladi (2022), Guo and Shmaya (2023)

Roadmap

- Model
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Single-project environment

- The agent can propose at most one project

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- The agent can propose at most one project
- We let $\alpha(u, v) \in [0, 1]$ denote the approval probability if the agent proposes project (u, v) , instead of using $\rho((u, v) | \{(u, v)\})$

Intuition: deterministic mechanisms for now

- Suppose that only deterministic mechanisms are allowed

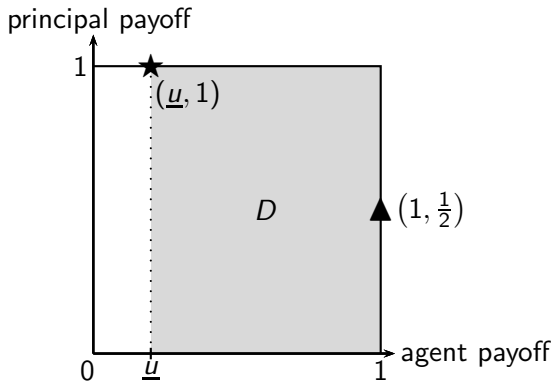
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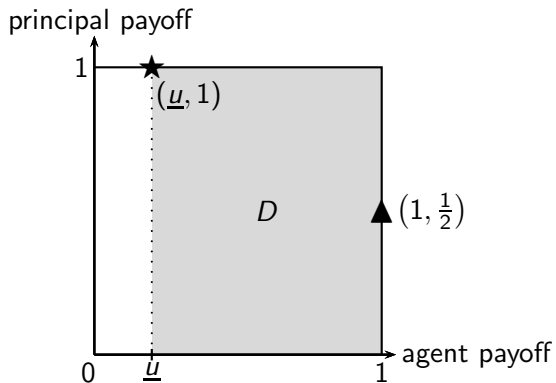
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- Shall the principal approve \blacktriangle ?

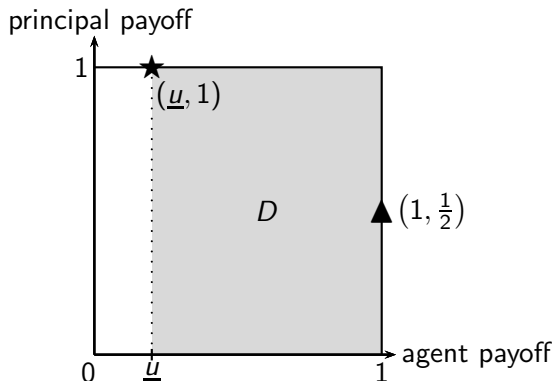


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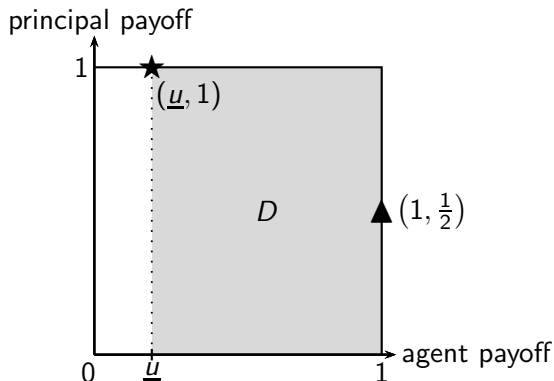
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Intuition: deterministic mechanisms for now

- If the principal approves \blacktriangle and $A = \{\blacktriangle, \star\}$, the agent will propose \blacktriangle and the principal suffers regret of $1/2$
- If the principal rejects \blacktriangle and $A = \{\blacktriangle\}$, the principal suffers regret of $1/2$



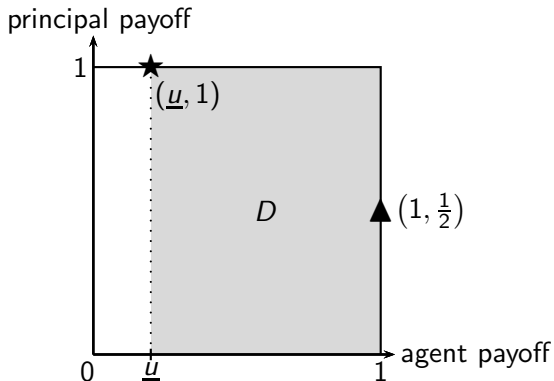
Deterministic mechanisms

Claim

In the single-project environment, the principal's worst-case regret under any deterministic mechanism is at least $1/2$.

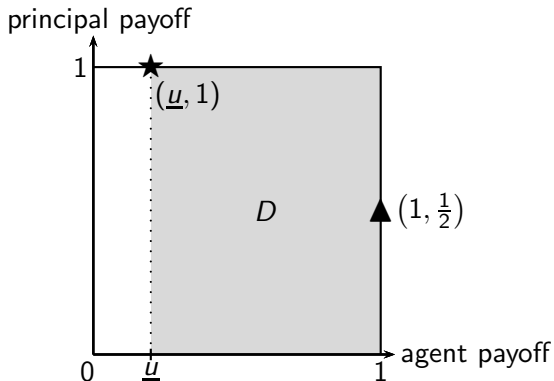
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- With randomized mechanisms, the principal approves \blacktriangle with probability \underline{u}



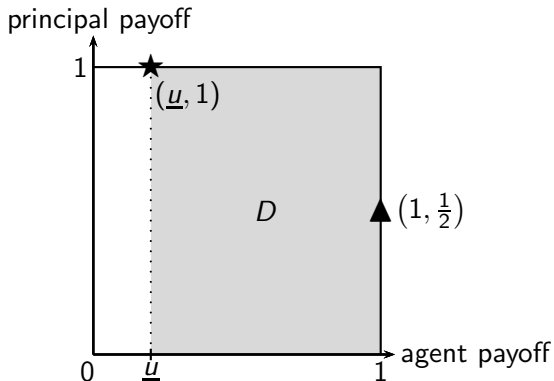
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- If $A = \{\blacktriangle\}$, the amount of inefficient rejection is reduced



Lower bound on the worst-case regret

Theorem

(i) The worst-case regret under any mechanism is at least R^s :

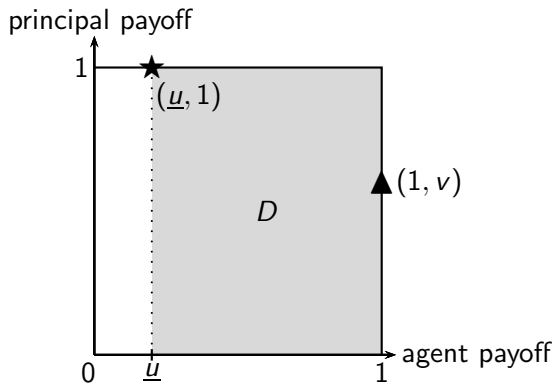
$$R^s \equiv \max_{v \in [0,1]} \min \{1 - v, (1 - \underline{u})v\} = \frac{1 - \underline{u}}{2 - \underline{u}}.$$

(ii) The two-tier mechanism:

$$\alpha^s(u, v) = \begin{cases} 1, & \text{if } v \geq 1 - R^s \text{ or } u = 0 \\ \underline{u}/u, & \text{if } v < 1 - R^s \text{ and } u > 0 \end{cases}$$

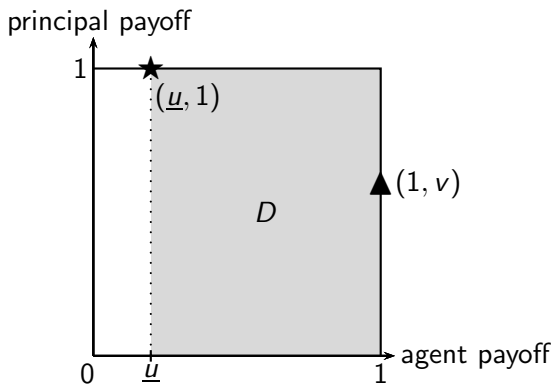
has the worst-case regret of R^s .

Proof: worst-case regret under any α is at least R^s



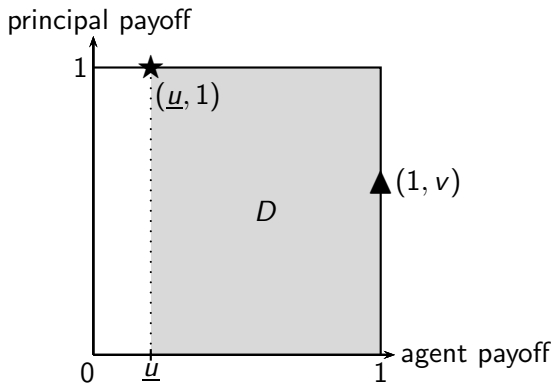
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- If $\alpha(1, v) > \underline{u}$, then if $A = \{\blacktriangle, \star\}$, the agent will propose \blacktriangle so regret is at least $(1 - v)$



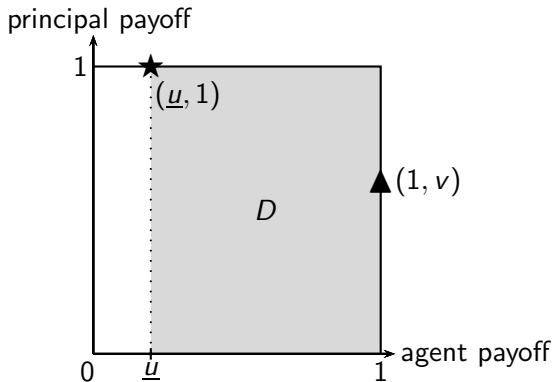
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 - If $\alpha(1, v) \leq \underline{u}$, then if $A = \{\blacktriangle\}$, regret is at least $(1 - \underline{u})v$
- \implies worst-case regret is at least $\min \{1 - v, (1 - \underline{u})v\}$ for any v



Optimal mechanism

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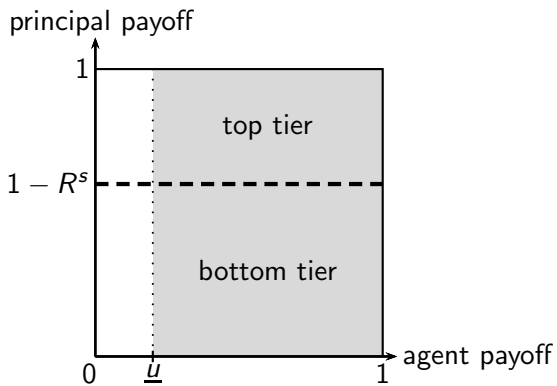
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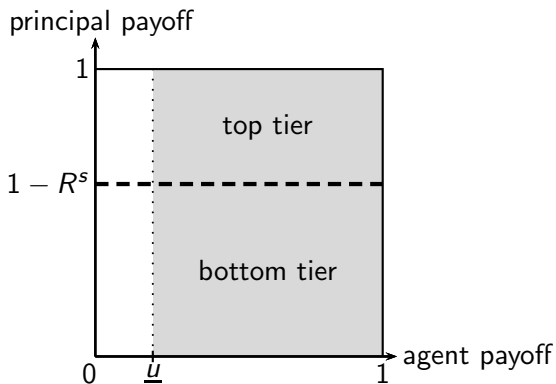
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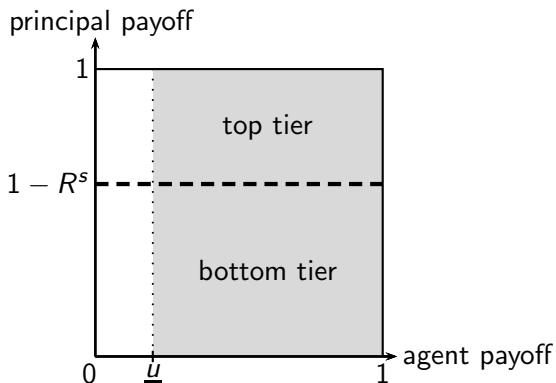
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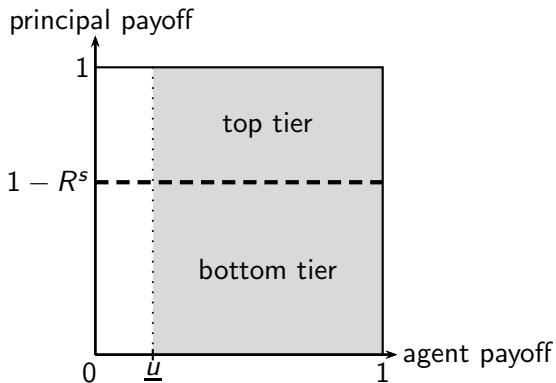


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- If he proposes a bottom-tier project (u, v) , his expected payoff is \underline{u}

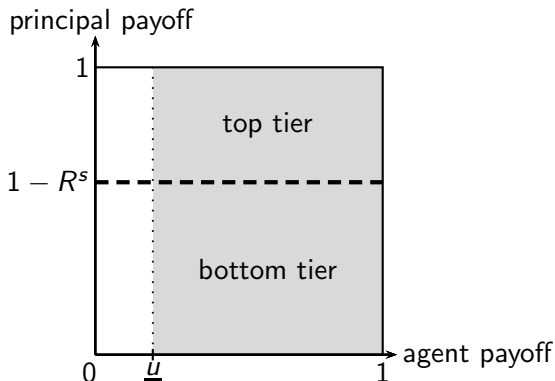


The agent's best response under α^S



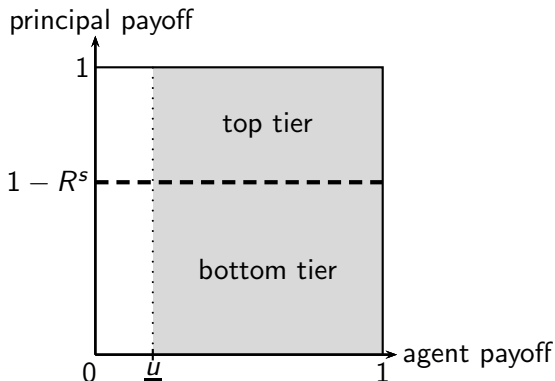
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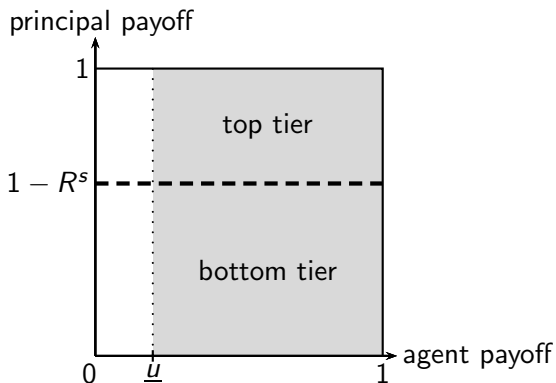


The agent's best response under α^s

- If the agent has projects with $v \geq 1 - R^s$, it is optimal to propose his favorite project among those with $v \geq 1 - R^s$
- Otherwise, it is optimal to propose a project that maximizes the principal's expected payoff $\alpha^s(u, v)v$

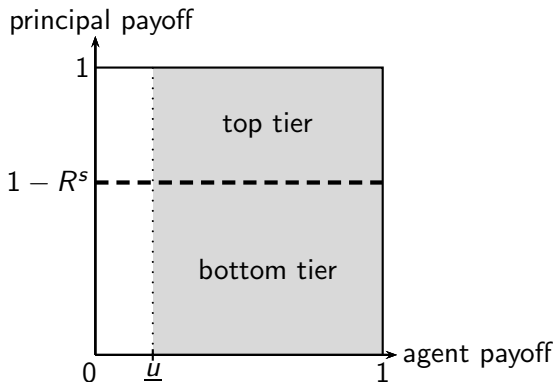


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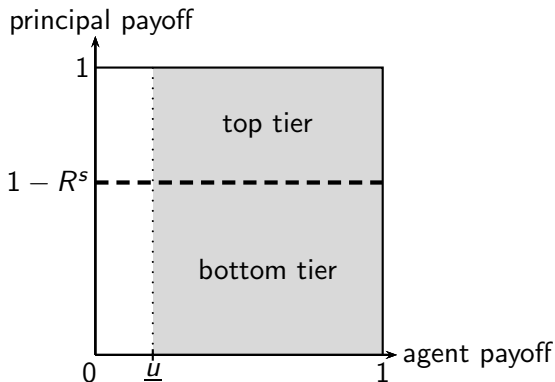
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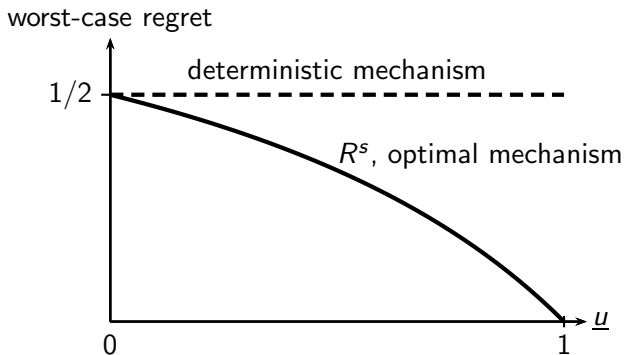
Proof: worst-case regret is R^s under α^s

- If the agent proposes a top-tier project (u, v) , it is approved for sure. Regret is at most $(1 - v) \leq R^s$.
- Suppose that the agent proposes a bottom-tier (u, v) . Let (u_p, v_p) be a principal's favorite project in A . Regret is:

$$v_p - \alpha^s(u, v)v \leq v_p - \alpha^s(u_p, v_p)v_p \leq (1 - \underline{u})v_p \leq R^s.$$



Worst-case regret R^s as a function of \underline{u}



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- The approval probability $\alpha^S(u, v)$ is monotone decreasing in u

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- The agent can propose any subset of the available projects, $P \subset A$

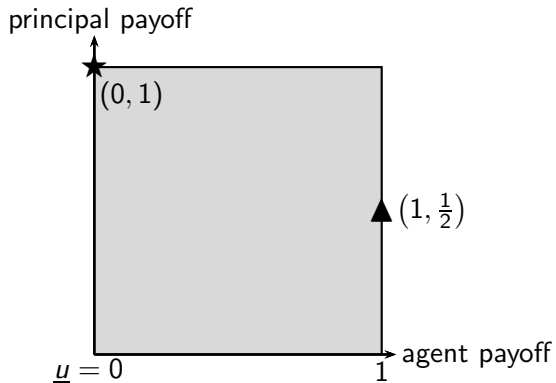
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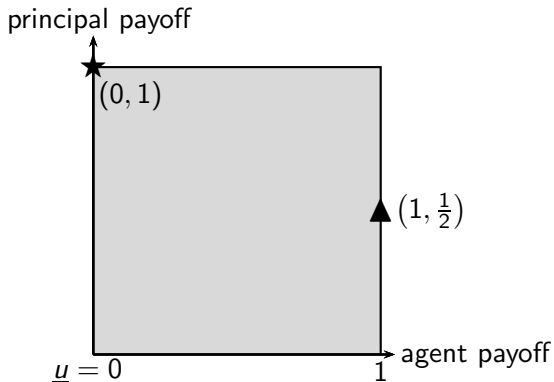
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- Given proposal P , $(u, v) \in P$ is chosen with probability $\rho((u, v)|P)$
- Revelation principle holds, so it is without loss to focus on mechanisms in which the agent optimally proposes all available projects, $P = A$

Intuition: single-project environment for now



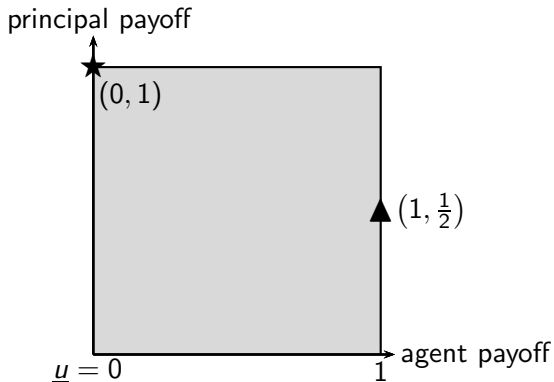
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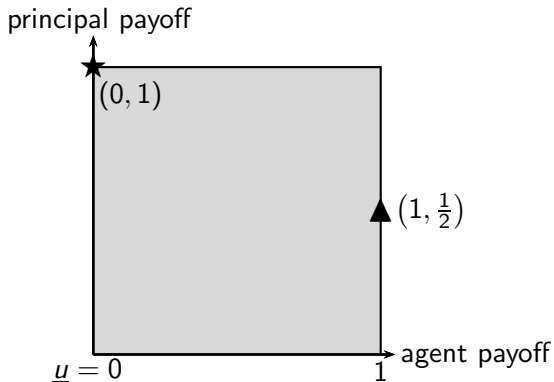
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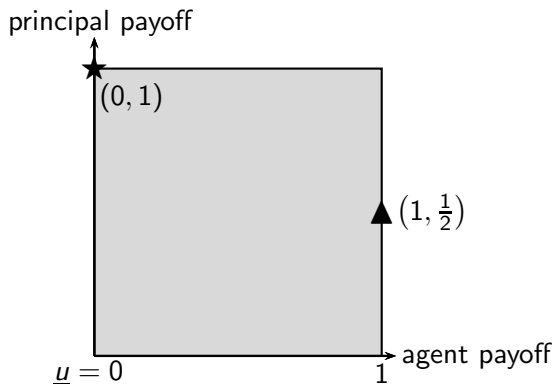


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- Worst-case regret is $R^s = 1/2$

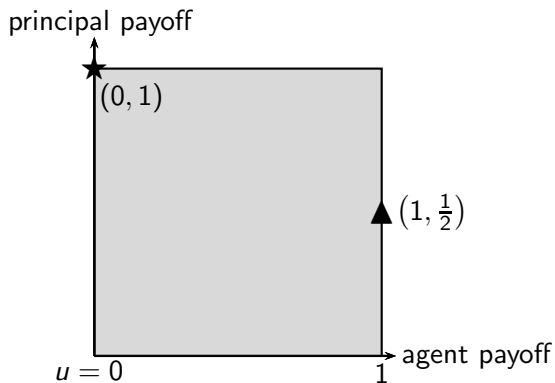


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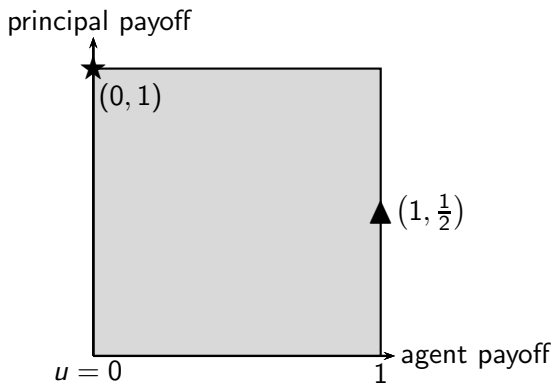
Intuition: how multiproject environment helps

- To incentive the agent to propose \star , his payoff from proposing $\{\blacktriangle, \star\} \geq$ his payoff from proposing \blacktriangle



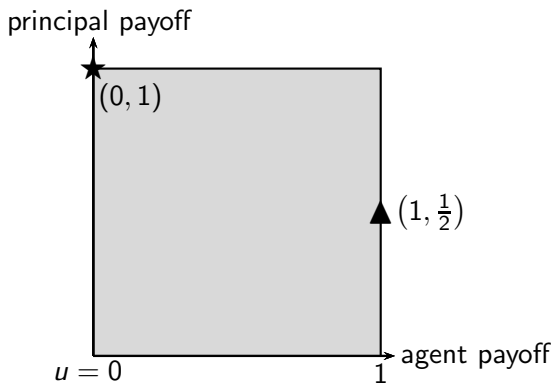
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- To incentive the agent to propose \star , his payoff from proposing $\{\blacktriangle, \star\} \geq$ his payoff from proposing \blacktriangle
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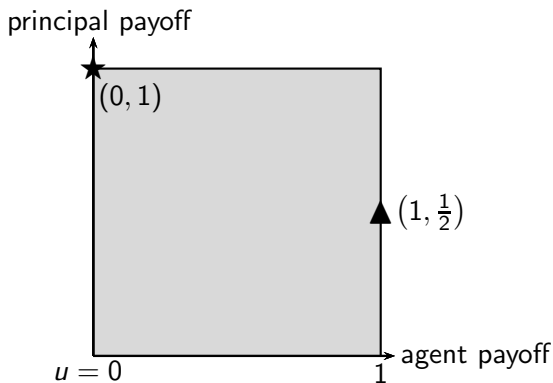
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Definition of PMP mechanism

Let $\alpha : [\underline{u}, 1] \times [0, 1] \rightarrow [0, 1]$ be a function.

The **proposal-wide maximal-payoff** mechanism (PMP mechanism) induced by α is as follows:

- If the agent proposes one project (u, v) , it is approved with probability $\alpha(u, v)$.
- If the agent proposes multiple projects, he is promised the maximal payoff from proposing each project alone (i.e., $\max_{(u,v) \in P} \alpha(u, v)u$).

Illustration of PMP mechanism

- If the agent proposes \blacktriangle alone, it is chosen with probability $1/2$ and the principal's payoff is y

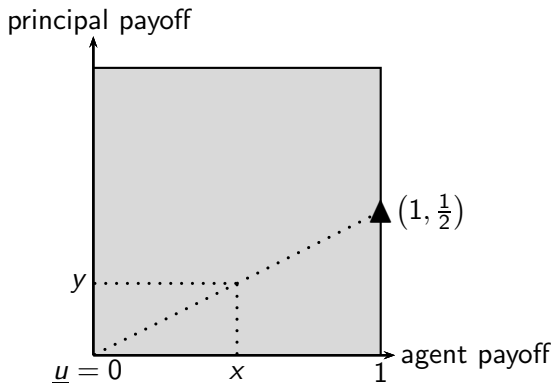
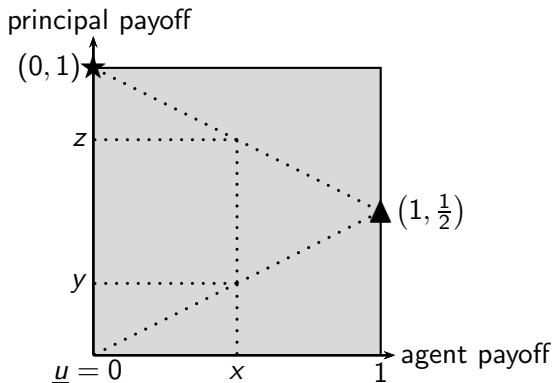


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- If the agent proposes \blacktriangle alone, it is chosen with probability $1/2$ and the principal's payoff is y
- The multiproject environment allows the agent to also propose \star , a better fallback option than rejection. The principal's payoff goes up to z



Lower bound on the worst-case regret

Theorem

For every $u \in [\underline{u}, 1]$ and $p \in [0, 1]$, let $\gamma(u, p)$ be

$$\gamma(u, p) = \min\{q \in [0, 1] : \underline{u} + q(u - \underline{u}) \geq pu\}.$$

(i) The worst-case regret under any mechanism is at least R^m :

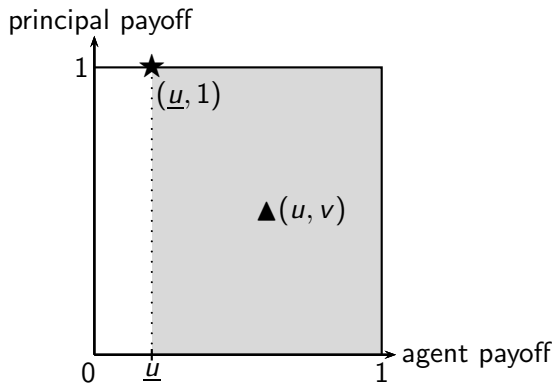
$$R^m = \max_{(u,v) \in D} \min_{p \in [0,1]} \max\{(1-p)v, \gamma(u, p)(1-v)\}.$$

(ii) Let ρ^m be the PMP mechanism induced by

$$\alpha^m(u, v) = \max\{p \in [0, 1] : \gamma(u, p)(1-v) \leq R^m\}.$$

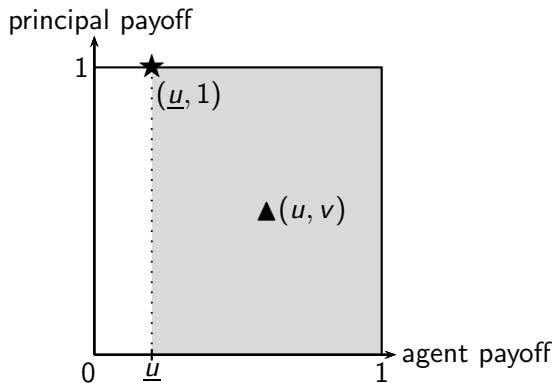
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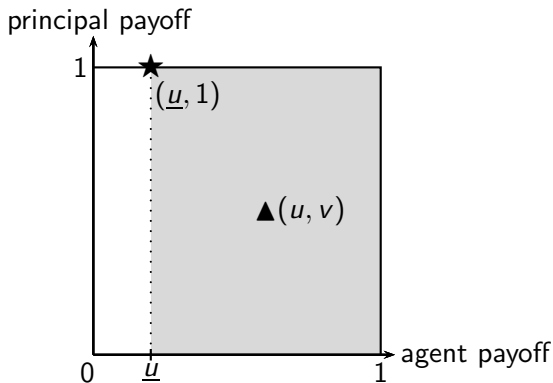
- Let ρ be the probability of approving (u, v) if it is proposed alone



Proof: worst-case regret under any ρ is at least R^m

- Let p be the probability of approving (u, v) if it is proposed alone
- If the agent proposes $\{(u, v), (\underline{u}, 1)\}$, the probability of choosing (u, v) is at least

$$\gamma(u, p) = \min\{q \in [0, 1] : qu + (1 - q)\underline{u} \geq pu\}$$

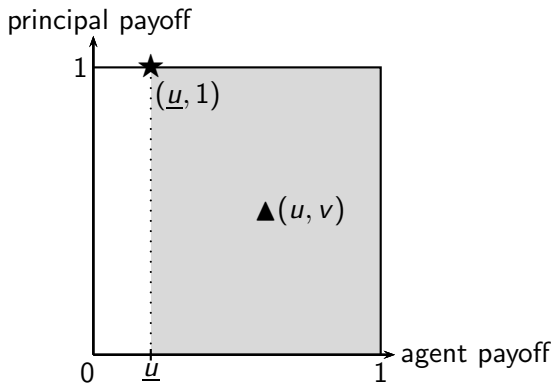


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Optimal mechanism

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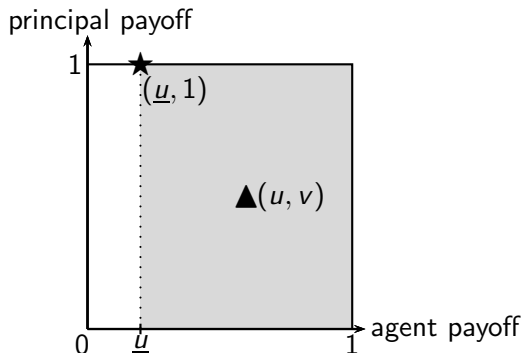
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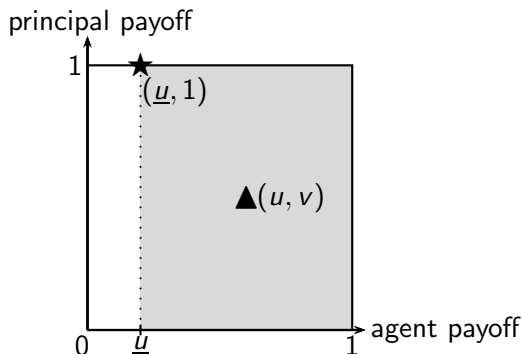
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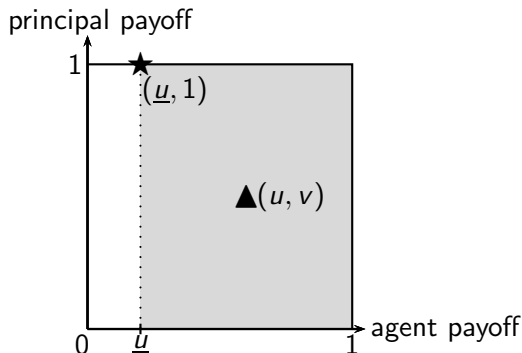
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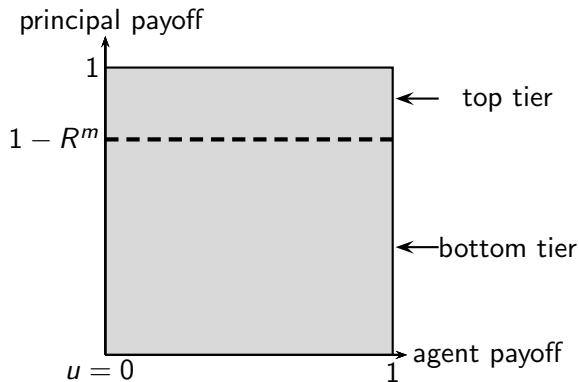
- $\alpha^m(u, v)$ is the highest probability p of approving (u, v) if it is proposed alone such that:
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 - and the regret from $\{(u, v), (\underline{u}, 1)\}$ is at most R^m



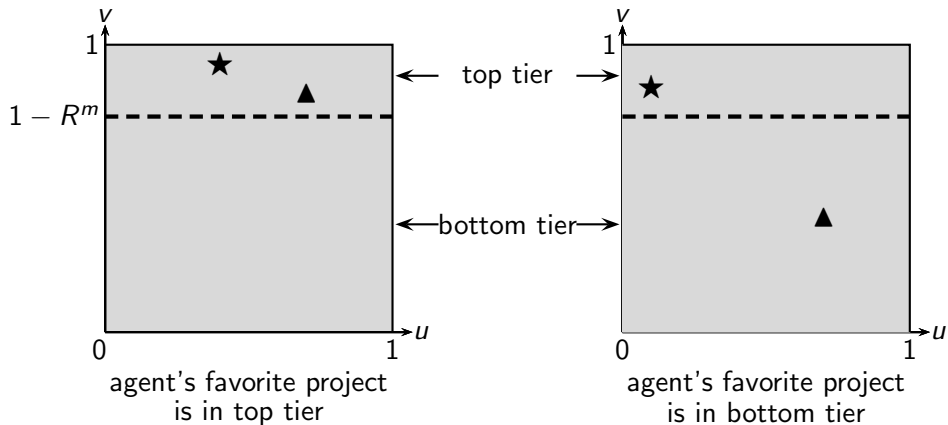
Two tier if the agent proposes one project

- The explicit expression for $\alpha^m(u, v)$ is:

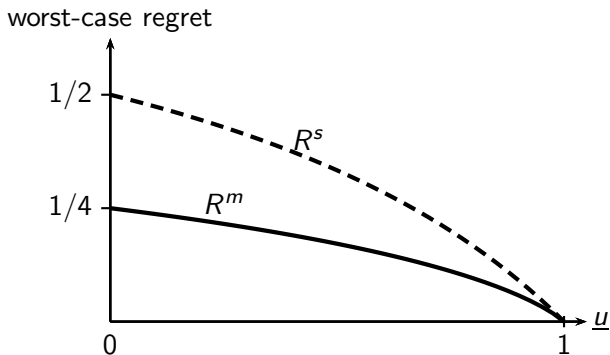
$$\alpha^m(u, v) = \begin{cases} 1, & \text{if } v \geq 1 - R^m \text{ or } u = \underline{u}; \\ \left(1 - \frac{R^m}{1-v}\right) \frac{\underline{u}}{u} + \frac{R^m}{1-v}, & \text{if } v < 1 - R^m \text{ and } u > \underline{u}. \end{cases}$$



When do several projects have a chance of being chosen



Worst-case regret: single-project vs multiproject



Roadmap

- Model
- Single-project environment
- Multiproject environment
- Discussion

Intermediate environments

- The agent can propose up to K projects

Proposition (Two is enough in minimizing worst-case regret)

For any $K \geq 2$,

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 - the project that maximizes $\alpha^m(u, v)u$
 - and the principal's favorite project;the corresponding choice function has the WCR of R^m .

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- The agent may need to exert effort to discover projects

Thank you!